ON THE USE OF A MACRO PROCESSOR WITH SUMX

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## I. INTRODUCTION

If the SUMX ${ }^{1}$ control statements are considered as forming a base "language", then the power of equipping it with a macro facility becomes obvious. In the following we describe a simple scheme which consists of making one preprocessing pass through a general purpose (that is, base-language-independent) macro processor called MACROS. ${ }^{2}$ The input text is prepared using base language statements (preferably in terms of variable symbols rather than fixed values) interspersed with preprocessor statements, such as macro definitions, macro calls, and value assignments to symbols. Prior to processing, this text is "compiled" by MACROS into a target text consisting entirely of base language statements.

For readers who are not familiar with, for example, macro-assemblers we list some of the advantages thus gained:
(1) Program parameterization.

This is achieved explicitly through the use of symbols for such quantities as number of channels, channel width, etc., the value assignments for which are delayed until the run time so that it is trivial to change them to any desired values. Furthermore names (symbols) are more convenient than numbers (BOUT locations).
(2) Shorthand notation and repetitive text generation.

This aspect of a macro processor is like SUBROUTINE in FORTRAN, apart from in- vs. out-of-line distinction. Many base language statements can be compressed into one macro call, if necessary, with variable arguments.
(3) Library facility.

Definitions for general purpose or frequently occurring macros may be collected into an external library to be shared among different jobs.

As discussed in Ref. 3 the advantages listed above are common features of macro processors. There are further advantages peculiar to our application as described below.

For experiments with large statistics it is desirable to minimize the "length" of the input data set (Data Summary Tape). One solution is to have for each event on a DST only the essential physics information such as the energy-momentum four-vectors and not the derivable quantities like invariant
masses, momentum-transfers, and decay angles. They are then calculated by means of CHARMs during the SUMX run. A peculiar advantage arises in our scheme because one and the same macro call is used to generate names for quantities of interest, prepare the corresponding CHARM calls to calculate them, and assign BOUT locations for storing the named results. The chance for error is thus greatly reduced and since each run is explicitly selfcontained it is easy to cross-check.

The present scheme was developed by John Ahern and the author as members of experimental group D at the Stanford Linear Accelerator Center and has been in use at SLAC since 1968 .

As described elsewhere ${ }^{2}$ MACROS is written in PL/I taking advantage of its list-processing facility with a small part dealing with the operating system such as that for requesting core-space in assembler language. All CHARMs and SUMX related routines are written in FORTRAN IV except those "pots and pans" ones such as vector- and matrix-manipulating routines which are coded in assembler language. These programs are running on the IBM system $360 / 91$ as implemented at SLAC and with some minor differences on a similar system at Max Planck Institut fur Physik und Astrophysik.

The capability of the present version of the macro processor, MACRO01, is somewhat limited because it lacks macro-time features such as macro-time variables (to do arithmetic with) and conditional branches (to be able to define macros recursively).

In Section II our scheme is described by an actual example with enough variety to illustrate most features of MACROS. Appendix A gives a formal description of the processor for reference purposes. In Appendix B we briefly summarize various CHARMs used in the example. In Section III several examples of errors in using the processor are collected, which also serve to show the limited capability of MACRO01. Appendix C has models for Job Control Language statements required to run.

## II. AN EXAMPLE

We illustrate the use of MACROS in SUMX by the example of Table I (referred to as A), which consists mostly of preprocessor statements (PPS) mixed with SUMX control (base language) statements. After one pass through MACROS this text turns into the target text consisting entirely of the base language statements as shown in Table $I I$ (referred to as B), which is then used to control SUMX. The macro library used in this example is shown in Table III (referred to as L).

Each PPS is one card long (columns 2-80) and is identified by the warning marker \% in the first column. The use of the warning marker serves to speed up the processing time. Comments can be added after the marker ; as in line 1A. Line 2A calls for a macro TWINKLE. Since it is not defined in the text scanned so far the macro library is searched. MACROS being a onepass processor macros must be defined and values must be assigned to symbols before they are used. This has the advantage of allowing local redefinitions and reassignments.

Lines 96-102L for TWINKLE show how to define a macro. Symbols (escape names) to be substituted are preceded by the escape character \& , double escape characters \&\& meaning the value must be placed starting at the specified column (left-justified). In MACROS no distinction is made between global (inter-macro) and local (within macro) symbols; any of \&A, $\& B$, or $\& C$ in TWINKLE may be left out of the parameter list, in which case values can be assigned by the assignment $\Leftrightarrow$ statement appearing anywhere before use.

A macro definition is terminated by MEND, except when it is followed by another macro definition. In such case the missing MEND is assumed by the processor, since the present version does not allow a nested macro definition.

Returning to TWINKLE the parameter A can be a character string like '1000, 100 ' or 'KEEP 2'. Unsigned integers and fixed point numbers (digits with a decimal point) do not need be surrounded by ${ }^{\top}$ as in the second argument in TWINKLE. Line 1A expands into lines 1-6B.

Line 3A calls for a macro LITTLE which, as defined in lines 103-108L, illustrates nested macro calls. Other macro calls up to line 7A are similarly straightforward. As seen in the expansions (lines 7-62B) the main purpose

## TABLE I

## INPUT TEXT

```
1 %; BEGIN 3 PRONG SUMX DECK
2 % CALL TWINKLE('1000,200',1,'PHOTON+P ... P + 2 PIONS 4/72')
3% CALL LITTLE(.005)
4 % CALL T33
5 % CALL MT3('AND',1,14,12,4,'PION CHANELS')
* CALL STAR
% CALL MT4('AND',1,14,12,4,29,'E>5.5')
* CALL C 33
*SELECT
* CALL BET(40,M45,.62,.86,'RHO')
*BLOCK6
* CALL SYMBDL
% WG T=10
14 % MACRO ANGLE (C,CS,PH)
15% CALL COS(C,CS)
1 6
17%
18 EVA
1 9
20
20
22
3. NPT=40
CALL DELSQ('RHO',045)
24% CALL ANGLE('RHO IN T-CHAN HEL FRAME',CJ45,PJ45)
25% CALL ANGLE('RHO IN S-CHAN HEL FRAME',CH45,PH45)
26
26 * NPT='!
2 7 ~ F I N I S H ~ 1 , ~
28 % MEND
29 % CALL WONDER(26)
30 % CALL WONDER(25)
31 % CALL WONDER(24)
32% CALL WONDER(2 3)
33 *ALL DONE
```

TABLE II
OUTPUT TEXT


Table II (continued)


## MACRO LIBRARY



Table III (continued)


Table III (continued)

```
131 % D45=731
132 % A 34=734 ; DECAY ANGLES
133 & A 35=744
134 % A45=754
135 % CJ45=756 ; J (H) REFERS TO T-CHANNEL (S-CHANNEL) HELICITY FRAMES
136 % PJ45=757 ; C (P) FOR COS (PHI)
137 % CH45=758
138 % PH45=759
139 % CW45=760
140 % PW45=761
141 % N=2
142 *CHARM
143 % CALL MCHM (34)
144 % CALL MCHM(35)
145 % CALL MCHM(45)
146 COSP- 0
147 cosP+ 0
148 COS+- 0
149 % CALL ACHM (45,12)
150 % CALL ACHM (34,21)
151% CALL ACHM (35,21)
152 %; MACROS TO SELECT CHANNELS
153 % MACRO H2 (N1,N2,N3,11,12,D)
154 % CALL EQU2(N1,'-4',I1,12,'8D.UNIQUE')
155 % CALL EQU2(N2,'-4','-8I1','-8I2','8D.AMB')
156% CALL MT2('',N3,N1,N2,'8D.ALL')
157 % MACRO T33
158 % CALL H2 (2,3,4,301,302,'P PI+ PI-')
159 % CALL H2 (6,7,8,303,304,'P K+K-1)
160 % CALL EQU2(9,'-4',305,'-305','P PBAR P')
```

of these is to define the master test (TEST 1) to eliminate unwanted events before CHARM calls are made to calculate various quantities to be SUMXed. The latter is done by a call (line 8A) to C33 (lines 122-151L). As shown C33 in particular assigns BOU' locations to various names for quantities being calculated.

MCHM (lines 116-117L) shows how a symbol can be concatenated and then substituted. Note here that the variable N is not included in the parameter list. Since it is comparatively slowly varying it is set by assignment (line 141L) rather than explicitly as an argument of the macro. C33 expands into lines 63-72B. For completeness explanations for CHARMs used are given in Appendix B.

Unassigned escape names, i.e., those which are preceded by \& but have not yet appeared in assignment statements are left unchanged in the text. This choice rather than automatic null-assignment is made because of possible use in other language such as in FORTRAN IV in which the character \& is used to designate statement labels. The purpose of SYMBOL (lines $1-31 \mathrm{~L}$ ) is to assign default values to symbols occurring globally as in the BLOCK6 and 7 macros (lines $33-50 \mathrm{~L}$ ), any of which may be set and reset before use as in lines 13 , 22 , and 26A.

Macros can, of course, be defined outside the library by the user as in lines $14-28 \mathrm{~A}$. The corresponding calls (lines $29-32 \mathrm{~A}$ ) result in SUMX control statements for a set of histograms repeated for different incident energy intervals as shown in lines 77-180B.

It was often asked, "in this scheme how many lines can be produced by a single line of typing?". The answer is obviously, "one-to-all", because all of the lines say, in this example can be lumped into a single macro. The need for doing so may conceivably arise in on-line applications.

It is also said that some of these features are available through additional programming in SUMX. This must be self-evident, because MACROS itself is a piece of program. The idea being propounded is in the use of a general purpose macro processor. In this respect it would have been better, if the "macro part" of the 360 assembler could easily be detached so that it could also be used elsewhere.
'Trivial syntax errors are evident from diagnostic messages from MACROS. In the following we describe examples of errors which require some explanation.
(1) Scanning for escape names is done from left to right. A peculiar effect can arise from concatenating names. For example in the use of the MACRO ACHM defined in lines 120-121 of Table III:

```
A12=1234
```


which is as expected. However,
$A=1$ '
CALL ACHM(12,34)
4
12
88 N
12
34
a EE 1
(2) When a substitution for an escape name is done, the line of text is rescanned (because the value itself may be an escape name) until all possible substitutions are carried out. This can cause a loop as in the following example:

MACRO M1(A)
\&A STAR
$\begin{array}{ll}E & \text { \&MENO (GENERATED) } \\ E\end{array}$
MACRO M2(A)
CALL M1('\&A (ITTLE')
MEND
CALL M2 ('TWINKLE')
**ERROR IN STMT $\quad 231, L E V E L=2, M A C R O=M 1$
,MODEL STMT $=223, L A S T$ TOKEN='A'
SEV $=8$ SIG PART OF STY1/RFPL VALUE LOST IN PACKFD REPL
***ERROR IN STMT $231, L E V E L=2$,MACRO=MI ,MODEL STMT
SEV $=8$ SIG PART OF STMT/REPL VALUE LOST IN PACKEO REPL $\quad 8$ 2 $3, L A S T$ TOKEV='A
***ERROR IN STMT $231, L E V E L=2, M A C R O=M 1 \quad$, MODEL STM SEV $=8$ SIG PART OF STMT/REPL VALUE LOST IN FACKFO REPL
** * ERROR IN STMT $231, L E V E L=2, M A C P O=M 1 \quad, M O D E L$ STMT $=$
$\begin{array}{ll}\text { RROR IN STMT } 231, L E V E L=2, M A C P O=M l \\ \text { SEV }=8 & \text { SIG PART DF STMT/REPL VALUE LOST IN PACKED REPL }\end{array}$
** $\ddagger$ ERROR IN STMT 231 ,LEVEL = 2,MACRO=M1 , MODEL STNT $=$
SEV = 8 SIG PART OF STYT/REPL VALUE LOST IN PACKED REPL
***ERROR IN STMT $231, L E V E L=2, M A C R O=M 1 \quad, M O D E L$ STMT =
SEV $=8$ SIG PART OF STMI/RFPL VALUE LOST IN PACKFD REPL
** $\#$ FRROR IN STMT $231, L E V E L=2, M A C R O=M 1 \quad$,MODEL STMT $=$ SEV $=8$ SIG PART OF STYT/RFPL VALUE LOST IN PACKED REPL
***ERROR IN STMT $231, L E V E L=2, M A C R O=M 1 \quad, M O J E L ~ S T M T=$ Z23,LAST TIKEN='AS




CALLED, IN STATEMENT 00789, FROM PROCEDURE WITH ENTRY POINT MACROS

| 299 | 0 | $\% ;$ |  |
| :--- | :--- | :--- | :--- |
| 300 | 0 | $\% ;$ |  |
| 301 | 0 | $\% ;$ |  |
| 302 | 0 | $\% ;$ |  |
| 303 | 0 | $\%$ |  |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |

## APPENDIX A

The following is reproduced from Ref. 2.

STMT LEVEL REPLC


SOURCB STATERENT HACO1 8SEP69

```
This section will describe the elenents of the nacros language in detail. The examples of preceding sections vere of an introductory nature, while this section is desigaed for reference purposes.
The syntax notation nses names of itens enclosed in < > syabols to denote syntactic entities. Definitions of then are either given in rords or in terns of other entities. This is indicated by an <iten> folloved by \(:==\) and definitions. The syabol \(\mid\) is used to indicate alterates.
```

```
I. ELEHEMTS
```

I. ELEHEMTS
<identifier> - 1-16 alphanumeric characters, the first of
which uust be alphabetic. Letters are A-Z,S,\&,a; digits
are 0-9. Identifiers are used to denote processor
variables and macros.
<integer> - 1-9 digits. Its usage in macros is identical to
strings, since this version has no macro-time arithmetic.
<string> -- there are tvo foras:
(1) Text string- 0-80 characters enclosed in quotes.
A quote vithin a string is represented by 2 consec-
ative quotes. The string of no length,or null string
is represented by ''.
(2) Fixed point number--string of digits preceded by, or
folloved by, or containing a decimal point. This
construct is included for convenience in writing
nacros statements; but it is not a number and does
not possess a numeric value. It should be thought
of as if it vere enclosed in quotes.
<constant> ::= <integer> | <string>
<formal paraneter> ::= <identifier>
The formal parameter is declared by appearance in a
MACRO statement. Its value is defined vhen the vacro
is expanded-rules are given belov.
<simple variable> ::= <identifier>
Simple variables are declared and assigned values
by appearance on the left side of an assignment statesent.
They assuse the type of the right side of the statement--
in this version of nacBOS it can alvays be considered
string type simce no macro~time arithmetic is permitted.
<processor variable> ::= <simple variable> | <foralal parameter>
<processor expression> ::= <constant> | <processor variable>
The processor erpression is used on the right side of
assignment statements or in actual parameters of CALL
statements. Its value is either the string value of
the constant or the value currently associated with
the processor variable. If the variable bas no value
associated vith it, then the expression is erroneous,

```
\begin{tabular}{|c|c|c|c|}
\hline 419 & 0 & \＄ & and its value is taken to be the noll string． \\
\hline 420 & 0 & \＄： & \\
\hline 421 & 0 & \％； & Examples： \\
\hline 422 & 0 & 8 ； & ＜identifier＞－－ABCD \＄\＄\＄OTHERI \\
\hline 423 & 0 & \％ & ＜integer＞－－ \(0 \quad 123123456789\) \\
\hline 424 & 0 & E； &  \\
\hline 425 & 0 & X； & ． 123 12．3 456．（Fixed point number） \\
\hline 426 & 0 & 8 ； & ＜processor expression＞－－ABCD 123 ．9＇STRING＇ \\
\hline 427 & 0 & 晏； & \\
\hline 428 & 0 & K； & \\
\hline 429 & 0 & 最； & IT．PROCESSOR S＊ATEAERTS \\
\hline 430 & 0 & \＄； & \\
\hline 437 & 0 & 5： & General： \\
\hline 432 & 0 & \＄i & Processor statements always have the character is\％in \\
\hline 433 & 0 & ． 8 & colum 1．The body of the statement ust be in columns 2－72． \\
\hline 434 & 0 & \％； & No continaation statesents are alloved．Dith the exception of \\
\hline 435 & 0 & 3： & assignment and nuil statesents，they begin vith a reserfed vard \\
\hline 436 & 0 & 8： & folloved by an operand field．Comments may follow the \\
\hline 437 & 0 & 景； & logical end of the statement vhen separated by a semi－colon \\
\hline 438 & 0 & \％； & （i）．Blanks may be used freely betyeen elenents and keywords， \\
\hline 439 & 0 & 8； & they are only reguired to separate alphanameric items． \\
\hline 440 & 0 & \＄i & Processor staterents are not subject to replacenent（section \\
\hline 441 & 0 & \％ & IV）or may they be generated by replacenent．Thus if a＂隹t is \\
\hline 442 & 0 & \％ & generated by replacenent in column t．nicBos will not cecognize \\
\hline 443 & 0 & \％ & the statement as a processor statement． \\
\hline 444 & 0 & 景： & Processor statenents are almays printed at top level and \\
\hline 445 & 0 & 景： & while a nacro is being edited．They are never printed at loyer \\
\hline 446 & 0 & 8： & levels（i．e．vithin a macro expansion）． \\
\hline 447 & 0 & \％； & \\
\hline 448 & 0 & E； & Mull or comment statement－ \\
\hline 449 & 0 & \％ & Fore：\(\quad\) ：comeents \\
\hline 450 & 0 & \％ & The null statesent is ignored by uncros．Then in a macro， \\
\hline 451 & 0 & 景： & it is not encoded，so no space is requiced for it in the \\
\hline 452 & 0 & 最； & editad Eacro． \\
\hline 453 & 0 & \％ & Example： \\
\hline 454 & 0 & \％ & X；The report you are reading is mostly null statements \\
\hline 455 & 0 & S； & \\
\hline 456 & 0 & \％ & MaCRO statement \\
\hline 457 & 0 & \％ & Pore：\(\quad\) AACBO＜macronale＞（＜formal parameter list＞） \\
\hline 458 & 0 & \％； & or \＆\(\quad\) HACRO＜nacroname＞ \\
\hline 459 & 0 & S； & ＜nacroname＞is：＜identifier＞ \\
\hline ＋60 & 0 & \(x_{i}\) & ＜foreal parameter list＞：： z （＜formal parameter＞ 1 \\
\hline 461 & 0 & \％ & ＜foraal parameter list＞，＜formal parameter＞ \\
\hline 462 & 0 & Ki & The MACRO statement heads a macro definition．The macror \\
\hline 463 & 0 & \％ & name is used in CaLL statements to reference the macro．It may \\
\hline 464 & 0 & 8 ； & be the same as processor variable identifier fo tyo macros \\
\hline 465 & 0 & \％i & may have the same name． \\
\hline 466 & 0 & 厚； & The formal parameter list，if supplied declares the \\
\hline 467 & 0 & 易： & identifiers as formal parameters．A maximu of ten may be given． \\
\hline 468 & 0 & 牙； & Details of their interpretation are in the next section． \\
\hline 469 & 0 & 第； & Macros say not be nested．The macro definition ends at \\
\hline 470 & 0 & \％ & a MEnD statement，the next maCRO statement，or end－of－file． \\
\hline 471 & 0 & 8 ； & Examples： \\
\hline 472 & 0 & 血； & 8 HaCRO FIBORACCI（I） \\
\hline 473 & 0 & \％ & \％Macro SoECLamations \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline STHT & LEVEL REPLC & & SOURCE STATEMEFT MACOI \\
\hline 474 & 0 & \％； & （ HACEO POACTION（TYP，ARG） \\
\hline 475 & 0 & N： & \\
\hline 476 & 0 & 7； & KEND statenent \\
\hline 477 & 0 & 亩； & Porme \({ }^{\text {\％}}\)（ BND \\
\hline 478 & 0 & \％； & The \(\quad\) EMD statesent terainates a macro definition that \\
\hline 479 & 0 & 车： & has not been othervise ended． \\
\hline 480 & 0 & 最； & \\
\hline 481 & 0 & 7； & Call Statement \\
\hline 482 & 0 & X： & Porw：\＆CALL＜macroname＞（＜actal parameter list＞） \\
\hline 483 & 0 & 易； & or \(\%\) CALL＜nacroname＞ \\
\hline 484 & 0 & 成： & ＜actual paraideter list＞：am＜actual parameter＞｜ \\
\hline 485 & 0 & 1： & ＜actual parameter list＞，＜actual parameter＞ \\
\hline 486 & 0 & 8； & ＜actual parameter＞：：＝＜processor expression＞ \\
\hline 487 & 0 & \％： &  \\
\hline 488 & 0 & 8 ； & which tust have been previously defined or be available on file \\
\hline 489 & 0 & 景； & SYSLIB． \\
\hline 490 & 0 & 8： & The actual parameter list if present is used to assign \\
\hline 491 & 0 & \％ & values to the formal parameters．Details are given in the next \\
\hline 492 & 0 & \％ & section． \\
\hline 493 & 0 & \％； & Examples： \\
\hline 494 & 0 & \％： & \％CALL PIBONAACI（ 1 ） \\
\hline 495 & 0 & 5： & \％CALL SDECLARATIONS \\
\hline 496 & 0 & \％ &  \\
\hline 497 & 0 & \％： & \\
\hline 498 & 0 & 戒； & TIThE statement \\
\hline 499 & 0 & \％： & Fore：\％TITEE＜string＞ \\
\hline 500 & 0 & 易； & or 8 TITLE \\
\hline 501 & 0 & X； & The TITLE statement sets the page title to＜string＞if \\
\hline 502 & 0 & \％； & present and causes a page eject．onitted operand leaves the \\
\hline 503 & 0 & \％ & title unchanged and causes a page eject only．The statement \\
\hline 504 & 0 & \％： & is not processed vithin a macro definition－only vhen it is \\
\hline 505 & 0 & \％； & expanded．The TITLE statement is never printed． \\
\hline 506 & 0 & \％ & Example： \\
\hline 507 & 0 & \％ & \％TITLE＂III．HMCBOS Language Description＂ \\
\hline 508 & 0 & \％： & The above was used to get the curcent page title． \\
\hline 509 & 0 & \％ &  \\
\hline 510 & 0 & \％ & Assigneent statement \\
\hline 511 & 0 & \＄： & Form：\(\quad\) \％＜leftside＞\(=\)＜processor expression＞ \\
\hline 512 & 0 & \％ & ＜leftside＞：：\(=\)＜processor variable＞ \\
\hline 513 & 0 & 8： & The assignment statenent is used to set the value of the \\
\hline 514 & 0 & \％ & leftside to the value of the processor expression．If the \\
\hline 515 & 0 & 8： & Ieftside is a processor variableg then the staterant also \\
\hline 516 & 0 & 易； & declares the identifier as such if it has not been used \\
\hline 517 & 0 & X： & before．If it is a formal parameter，then the assignaent \\
\hline 518 & 0 & \％； & is only retained during this expansion of the macro．The \\
\hline 519 & 0 & \％ & actual parameter，if it vas a processor variable，vill not \\
\hline 520 & 0 & \％ & be changed by the assignment． \\
\hline 521 & 0 & \％ & \\
\hline 522 & 0 & 5； & ITI MACRO EXPAMSIOMS ． \\
\hline 523 & 0 & \％： & III．MACRO EXPANSIONS \\
\hline 524 & 0 & \％ & \\
\hline 525 & 0 & 凧； & When a macro is expanded using the CALI statement，the \\
\hline 526 & 0 & 易； & following actions take place： \\
\hline 527 & 0 & \％ & （1）If any forimal parameters vere declared，any falues currently \\
\hline 528 & 0 & \％ & assigned to the formal parameter identifiers are saped on a \\
\hline
\end{tabular}

(2) pushdovn stack.

The values of the actual parameters supplied in the CALL statenent are assigned to the formal parameter identifiers. If amy or all of the actual parameters are onitted, then the corresponding formal parameters are set to the null string. Excess actual parameters are ignored. Thus a call statement of the form ( CALL FURCTIOM ( 3 )
for the nacro powction above, vill set TYP=" and aRG=" \({ }^{\prime \prime}\).
(3) nacros begins fetching input from the body of the macro definition. The processor variables that are currently active vill be used during processing--this includes any variables set by acros that called this one.
(4) When the end of the macro definition is reached, the values previousiy associated vith the formal parameter identifiers are restored vith the values saved on the pushdovn stack, if any. macros then resures fetching statements in the environnent of the invoking call statenent: but changes made to variables global to the called nacro are retained.
The earimun depth of macro calis is 10 levels.

\section*{If. REPLACERENT IR BASE LAMGAGE STATEBEATS}

Every base language statenent (i.e. no 'st in col 1) outside of a acro definition is scanned by hacaos for the escape character ' \(E^{\prime}\). which signals a potential replacement. If the ' \(\varepsilon\) ' is folloved by an identifier or by another ' \(\varepsilon\) ' and an identifier, and if the identifier is that of a processor variable that is currently active, then a replacenent is performed. The single anpersand is osed to denote "packed" replacement, for free forat applications; and the double appersand for "non-packed" replacenent, for use uhere colinn positions are critical.

The tern "escape name" is used to refer to the escape character (s) and the innediately following ldentifier. for one case of packed replacenent, a period deliniter is included in the escape nate.

Scanaing for an escape character proceeds fron left to right. The scanning field is the entire fortran statement, including contination lines, if the fortri option is used; else it is the colunns between the begin and end columas finclusive) of the 80 character record. The default scanning colums are 1 thru 80 ; they may be changed by the user if needed.

When an escape name if found, and replacement is done (the identifier has a value), then the scan is restarted from the begin colunn. If the identifier does not have a value, then the scan proceeds to the right. Rescanning continues until no
more escape characters renain, or no more replacenents can be
done. The rescanning mechanisi allous construction of escape names by replacenent. The maxinut number of replacements alloved in a base language statement is 30 . specific rales for the replacesent modes are as follovs:
\begin{tabular}{|c|c|c|}
\hline 584 & 0 & 景； \\
\hline 585 & 0 & 8； \\
\hline 586 & 0 & 置： \\
\hline 587 & 0 & \％； \\
\hline 588 & 0 & \％ \\
\hline 589 & 0 & 景： \\
\hline 590 & 0 & \(\boldsymbol{8}\) ； \\
\hline 591 & 0 & \％： \\
\hline 592 & 0 & \％： \\
\hline 593 & 0 & 宜； \\
\hline 594 & 0 & \％： \\
\hline 595 & 0 & 县： \\
\hline 596 & 0 & \％； \\
\hline 597 & 0 & \＄： \\
\hline 598 & 0 & \％： \\
\hline 599 & 0 & 8； \\
\hline 600 & 0 & \％： \\
\hline 601 & 0 & 男： \\
\hline 602 & 0 & 栜： \\
\hline 603 & 0 & 敏； \\
\hline 604 & 0 & \％； \\
\hline 605 & 0 & \＄； \\
\hline 606 & 0 & \％ \\
\hline 607 & 0 & \％ \\
\hline 608 & 0 & \％ \\
\hline 609 & 0 & 易； \\
\hline 610 & 0 & \％ \\
\hline 511 & 0 & 8； \\
\hline 612 & 0 & \％： \\
\hline 613 & 0 & \％ \\
\hline
\end{tabular}
（1）Packed replacement－This is indicated by a single＂E＊
folloved by an identifier，deliaited by any special
character．If it is delinited by＇．＇，then the period will also be removed from the text－－this allovs concat－ enation of a replacenent value with a letter or digit． Replacement is done by renoving the escape character and the identifier（and period，if necessary）from the tert，and inserting the value associated with the iden－ tifler in its place．The remainder of the statenent to the right of the escape name，extending to the end of the scanning field，is shifted right or left as needed to accommadate the replacement value．Blanks will be added at the end of the statenent if the value is shorter than the escape name，or truncation will be preformed if the value is loager．In the latter case，loss of non－blank characters will be noted by nicros．
（2）Mon－packed replacenent－This is denoted by a double escape character＇\(\varepsilon \varepsilon^{\prime}\) followed by an identifier．The identifier is delinited by any special character（no special rule for period applies）．The replacement of the escape name is done vithout any shifting of the part of the statement to the right of the escape nane． If the replacement valne is shorter than the escape nane，then blanks are added to it．If the value is longer， then blank positions folloving the escape nane are used to hold the trailing non－blank characters of the value． If there is still not enough roon，even using the blank positions，and discounting trailing blanks of the value， then hacros notes the truncation of the value．


SOURCE STATEMENT
HACO1 BSEP69

Which causes base language statements containing escape characters to be isted before replacenent is done. Norally, they are listed after replaceneat only.
SIZS (10000) The amount of sgace (in bytes) to be allocated to storage of the values of variables. Should one find that he has run out (error message), he can rerun vith a larger value (up to 32767). Space could run ont if a large number of variables are all active at the same time.

When nacros is run, the first page of the listing gives the options used, and a nuaber sIzy* which is the amount of free space in bytes that MACROS thinks it has for storage of strings. macro definitions and control information. macros uill use all the free space in the region in wich it is run. Minimar region size is about 100 k bytes.

Errors in \(\operatorname{HACROS}\) input are flagged with a text message indicating the type of ercor,disposition, and location. Severities are associated vith the errors ranging from 4-16. The highest severity is passed back to os as a return code that can be tested in 3CL to suppress subsegrent job steps.

\section*{APPENDIX B}

\section*{CHARMS}

An event in our SUMX is represented by a FORTRAN array \(P(4, N)\) consisting of four-vectors, one for each participant of an N-particle reaction,
\[
1+2 \rightarrow 3+4+\ldots
\]
where numbers correspond to the second index of the array P. As discussed in the text the purpose of the CHARMs to be described here is to calculate from \(P\) various quantities of interest, which are to be specified at the SUMX run time.

As described in the SUMX manual \({ }^{1}\) a CHARM control-card allows five parameters, L1, L2, ..., L5, by means of which user can communicate with the requested CHARM subroutine. In the descriptions below "pointer" means BOUT location, a group of particles is denoted by ij. .., the total number in this group by N. The CHARM parameter L5 is always the base pointer for the array \(P\).
(1) CHARM2

Invariant mass and four-momentum transfer for a group of particles.
L1: pointer for answer
L2: N
L3: ij...
\(\mathrm{L} 4: 0\) to get \(\mathrm{M}=\mathrm{SQRT}\left[\left(\mathrm{P}_{\mathrm{i}}+\mathrm{P}_{\mathrm{j}}+\ldots\right)^{2}\right]\),
-n to get \(\Delta^{2}=-\left(\mathrm{P}_{\mathrm{i}}+\mathrm{P}_{\mathrm{j}}+\ldots-\mathrm{P}_{\mathrm{n}}\right)^{2}\).
(2) CHARM3

Four-momentum transfer \(\left(\Delta^{2}\right)\), production cosine (COS) and momentum (Q) for a group of particles in the overall center-of-mass.

L1: base pointer for answer vector A
\(\mathrm{A}(1)=\mathrm{COS}\),
\(A(2)=\Delta^{2}-\Delta_{\text {min }}^{2}\) (kinematic minimum),
\(A(3)=Q\) (calculated only if \(L 1\) is negative).
L2: N
L3: ij...
L4: 1 or 2 to specify with respect to beam or target.

\section*{(3) CHARM4}

Two or three body decay angles.
L1: base pointer to answer vector A
L2: N (2 or 3)
L3: ij...
L4: ab ( 12 or 21 depending on where 1 or 2 is to be incident)
The calculated decay angles have the following meaning. Let
\[
R=P_{i}+P_{j}+\ldots
\]
and T be such that
\[
a+b=R+T
\]

Let \(\hat{y}\) be the normal to the production plane,
\(\hat{y} / / \overrightarrow{\mathrm{T}} \times \overrightarrow{\mathrm{a}}\) in the R-rest frame, or
\(/ / \vec{a} \times \vec{R}\) in the laboratory frame.
Then in the rest frame of \(R\) define the \(t\)-channel helicity (Gottfried-Jackson) axes as
\[
\hat{z}=\hat{a}
\]
and
\[
\hat{x}=\hat{y} \times \hat{z}
\]

The s-channel helicity axes are related to the above by a rotation about the common y-axis. Let
\[
\hat{\mathrm{z}}_{\mathrm{H}}=\hat{\mathrm{R}} \text { in the overall center-of-mass, }
\]
or
\[
=-\hat{\mathrm{T}} \text { in the } \mathrm{R}-\mathrm{rest} \text { frame }
\]
and
\[
\begin{aligned}
& \hat{\mathrm{y}}_{\mathrm{H}}=\hat{\mathrm{y}} \\
& \hat{\mathrm{x}}_{\mathrm{H}}=\hat{\mathrm{y}}_{\mathrm{H}} \times \hat{\mathrm{z}}_{\mathrm{H}}
\end{aligned}
\]

Then using \(i\) as the analyzer CHARM4 calculates the following
\[
\begin{aligned}
& \mathrm{A}(1)=\hat{\mathrm{z}}_{\mathrm{H}} \cdot \hat{\mathrm{z}} \\
& \mathrm{~A}(2)=\tan ^{-1}\left(\hat{\mathrm{z}}_{\mathrm{H}} \cdot \hat{\mathrm{x}} / \hat{z}_{\mathrm{H}} \cdot \hat{\mathrm{y}}\right)
\end{aligned}
\]
\[
\begin{aligned}
& \mathrm{A}(3)=\hat{\mathrm{i}} \cdot \hat{\mathrm{z}} \\
& \mathrm{~A}(4)=\tan ^{-1}(\hat{\mathrm{i}} \cdot \hat{\mathrm{x}} / \hat{\mathrm{i}} \cdot \hat{\mathrm{y}}), \\
& \mathrm{A}(5)=\hat{\mathrm{i}} \cdot \hat{\mathrm{z}}_{\mathrm{H}}, \\
& \mathrm{~A}(6)=\tan ^{-1}\left(\hat{\mathrm{i}} \cdot \hat{\mathrm{x}}_{\mathrm{H}} / \hat{\mathrm{i}} \cdot \hat{\mathrm{y}}_{\mathrm{H}}\right),
\end{aligned}
\]
where the angles are in degrees.
It appears in this scheme that because SUMX treats L3, L4 as integers particle index greater than 9 cannot be accommodated. One solution, as long as they need not be addressed at once, is to move base pointer \(L 5\) within \(P\) or to rearrange members of \(P\) before use. Far better solution is to "ask" SUMX to regard these parameters as character strings, so that a number system of arbitrary base can be employed.

\section*{APPENDIX C}

JCL EXAMPLES
We give examples of JCL statements necessary for runs on the System 360/91 as implemented at SLAC and at MPI. Catalogued procedures used are current ones in May 1972.

\section*{(1) At SLAC}
// JOB card
//JOBLIB DD DSN=WYL.ED.PUB. LIB9, DISP=(SHR, PASS)
//MACRO EXEC PGM=MACRO01
//SYSPRINT DD SYSOUT=A
//SYSOUT DD DSN=\&MOUT, DISP=(NEW, PASS), UNIT=SYSDA,
\(/ / \mathrm{SPACE}=(\mathrm{TRK},(200,10)), \mathrm{DCB}=(\mathrm{RECFM}=\mathrm{FB}, \mathrm{LRECL}=80, \mathrm{BLKSIZE}=3200)\)
//SYSLIB DD DSN=WYL.ED.JAP.SRC9(SMCR1), DISP=SHR
// DD user macro library
//SYSIN DD*
input text
//SUMX EXEC FORTHCLG, PARM. FORT='OPT=2'
//FORT.SYSIN DD*
user FORTRAN source, if any
//LKED. SYSLIB DD DSN=WYL.ED.PUB. LIB9, DISP=SHR
// DD DSN=SYS1. FORTLIB, DISP=SHR
// DD DSN=SYS3.FORTLIB, DISP=SHR
// DD DSN=SYS4. FORTLIB, DISP=SHR
//LKED.SYSIN DD*
INC LUDE SYSLIB(SUMX)
ENTRY MAIN
//GO.FT10F001 DD user DST description
//GO. SYSIN DD DSN=\&MOUT, DISP=(OLD, DELETE)
(2) At MPI
// JOB card
//D EXEC PGM=MACRO01,REGION=300K
//STEPLIB DD DSN=LOAD. JHP, DISP=SHR
\(/ /\) SYSLIB DD DSN=SOR.JHP(MACLIB), DISP=SHR
// DD user macro library
//SYSPRINT DD SYSOUT=A
//SYSOUT DD DSN=MOUT, DISP=(NEW, PASS), UNIT=DISK,
// SPACE=(TRK, \((200,10)\) ), DCB=SOR.JHP
//SYSIN DD*
input text
//S EXEC SUMX
//C.SYSIN DD*
user FORTRAN source, if any
//L. SYSLIB DD DSN=LOAD. JHP, DISP=SHR DD DSN \(=\) SYS1. FORTLIB, DISP=SHR
//G. FT10F001 DD user DST description
//G.SYSIN DD DSN=MOUT, DISP=(OLD, DELETE)

\section*{REFERENCES}
1. Although the technique is evidently general, we discuss a particular application to the CERN version of the SUMX described in J. Zoll, CERN Track Chamber Program Library Manual (1970).
2. John Ahern, MACROS-Statement Oriented Macro Processor, SLAC Computation Group User Note 29 (1969).
3. P. J. Borwn, A Survey of Macro Processors, Annual Review in Automatic Programming, Vol. 6, Part 2 (Pergamon Press, New York, 1969).```

